



# Cellule: Lightweight Execution Environment for Virtualized Accelerators

Georgia Tech

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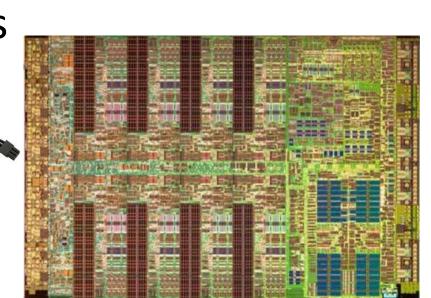
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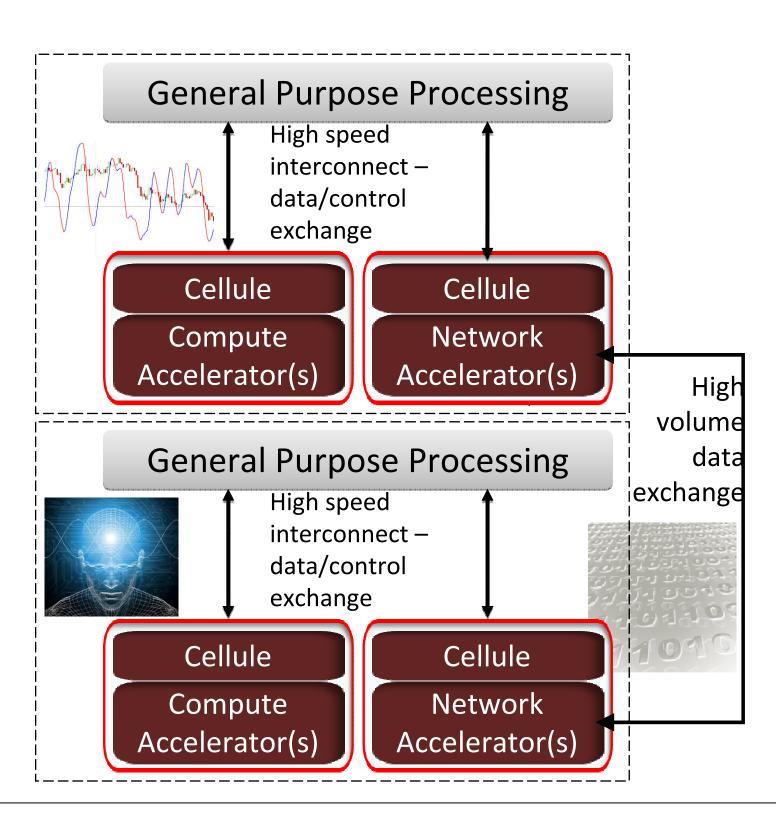
#### What is Cellule?



- Cellule small, custom execution environment for accelerators (cellule n., a small cell)
- specialized VM for access to and management of accelerator resources (e.g., STI Cell, Intel IXP, etc.).
- helps improve portability and attain high





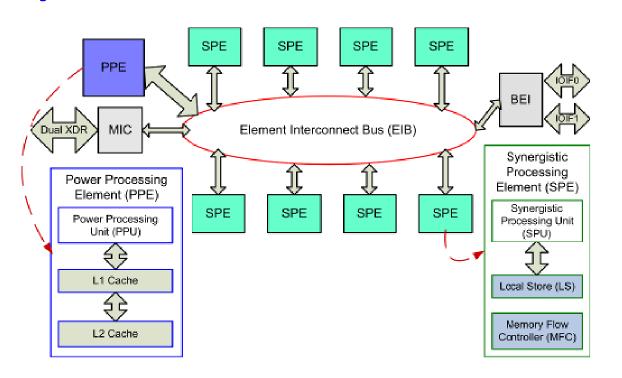


## Objectives

Improved performance and lightweight, specialized environment for accelerators

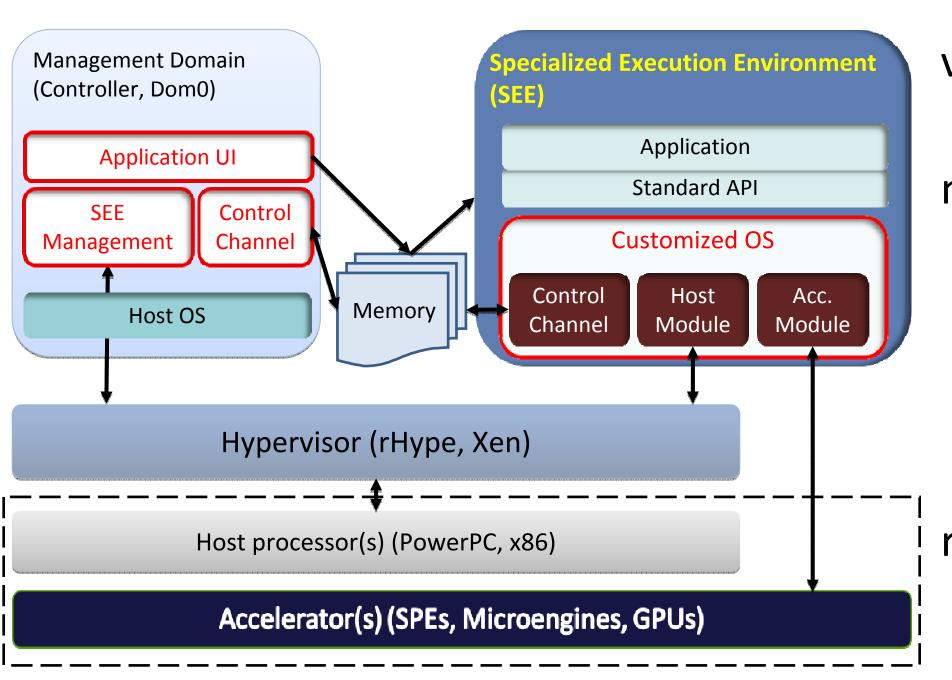
- Accelerator specific memory and execution model
- Simple task model
- Ease of customization and debugging
- Improved signaling mechanisms
- Finer grained scheduling and resource management

## Example Accelerator – Cell B.E



- Importance of asynchronous operation and PPE decoupling
- Architectural support for virtualization
- Lightweight environment conducive to high performance and customization

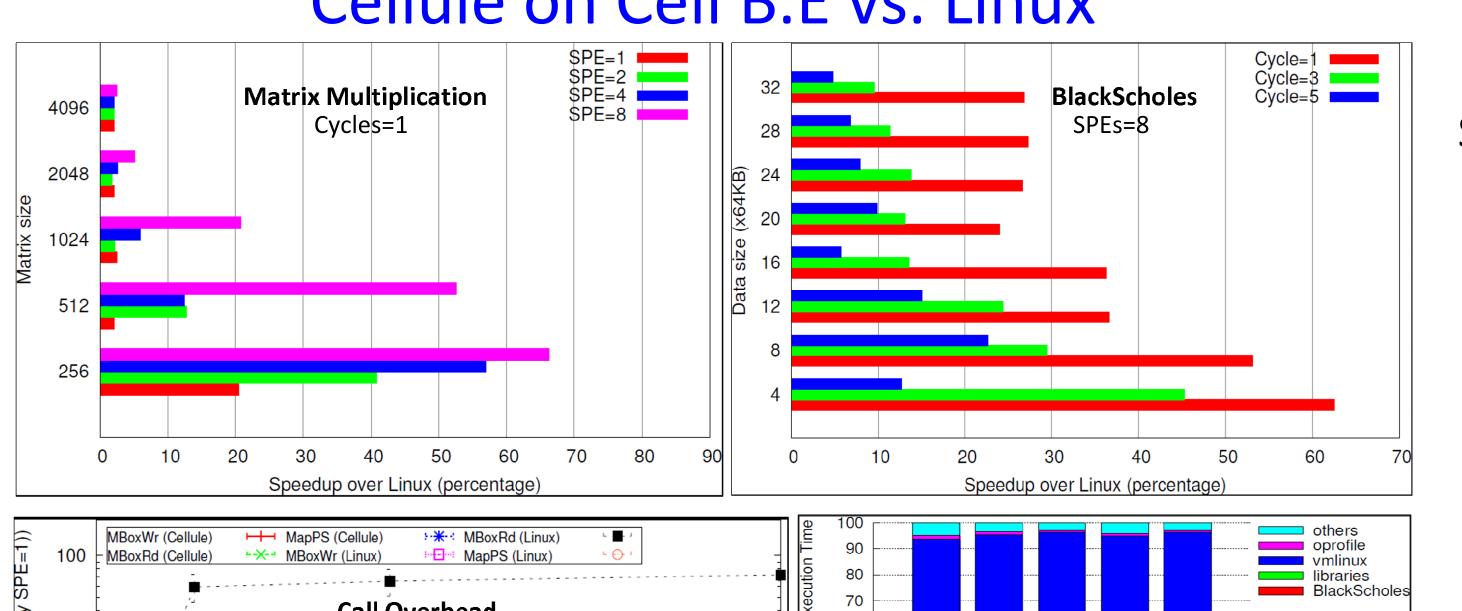
### Cellule Architecture

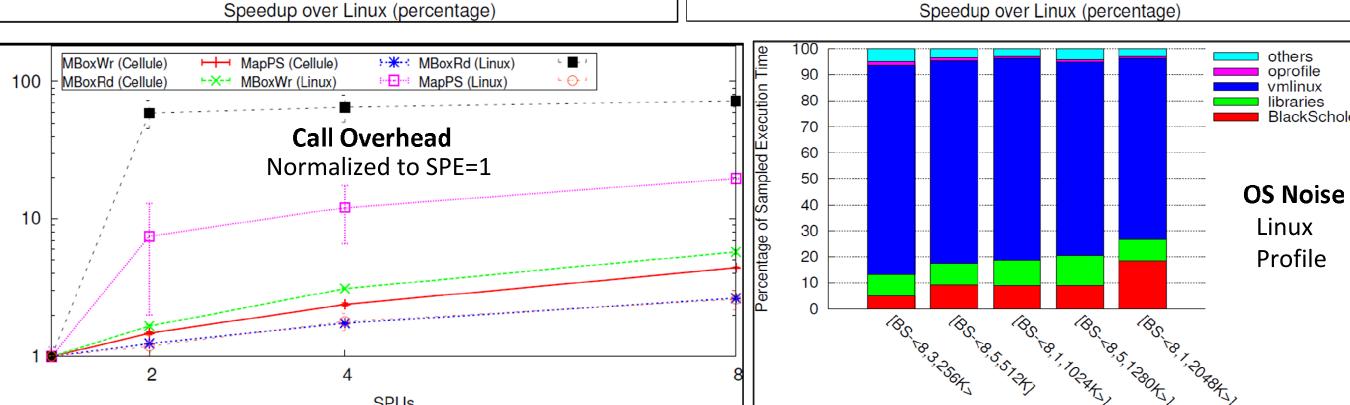


#### Discussion

- Adapt the architecture to various accelerators — Cell, IXP
- Use hypervisor and a management domain to
  - Provide familiar UI to user
  - Run the requested application in a lightweight 'special execution environment' (SEE)
- Hypervisor manages resources, multiplexes SEEs
  - Bypass it once accelerator acquired by SEE – fast access

#### Cellule on Cell B.E vs. Linux





#### Benchmarks

- Matrix multiplication and Black Scholes - variable
  - Number of SPEs
  - Number of iterations
  - Data sizes
- Experiments on Cell Blade QS20

# Results - highlights

- Cellule better than Linux mostly
- Improved setup and execution time due to adapted memory and simplified task model
- Particularly important for smaller problem sizes, i.e., `difficult' accelerator applications

# Cellule on IXP – 'Open' vs. 'Closed' Accelerators

IXP used as an open accelerator for fine-grained resource management

- Offers gray-box view to management domain:
- Open access to IXP processing engines and memory through well-defined kernel or hypervisor extensions
- Lightweight signaling for fine-grain resource-control
- Scheduler co-ordination opportunities:
- IXP and Host schedulers interact to tune asynchronous accelerator invocations with Host process or VM execution.
- 1 App, 1 Egress Port, 1 Ueng for 920 1 App, 1 Egress decoupled UEs 880 2 Apps, 2 Egress 870 Ports, decoupled Rx,Tx UEs 2 Apps, 2 Egress ports, decoupled Rx, 2 Tx UEs Netperf Application
- Netperf TCP throughput increases by dedicating IXP resources to certain packet-flows
- Fine-gain control of IXP micro-engines to achieve throughput differentiation

## Ongoing Work

- On Cell:-
  - Optimization of SEE, rHype with thread, SMP support
- Sophisticated scheduling possibly with dynamic partitioning of SPEs
- IO over 9P communication channel to achieve RoadRunner configuration
- Extend Cellule to other accelerators
- On IXP:
- Expressive message abstractions between Host cores and Accelerator cores for better coordination

# References

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